

Speeding up splits in Hoeffding Tree Regressors

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- Incremental regression trees
- 2 The (incremental) variance problem
 - Naive algorithm
 - Welford's algorithm
- 3 Extended Binary Search Tree Observer
- 4 Quantization Observer
- **5** Benchmarking





Disclaimer

This presentation is grounded on the following paper:

Mastelini, S.M. and de Leon Ferreira, A.C.P., 2021. Using dynamical quantization to perform split attempts in online tree regressors. *Pattern Recognition Letters*, 145, pp.37-42.



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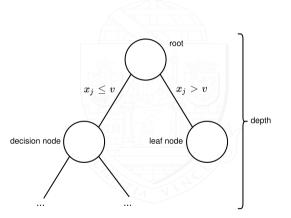


Introduction

Context:

- Numerical input features
- Axis-aligned splits

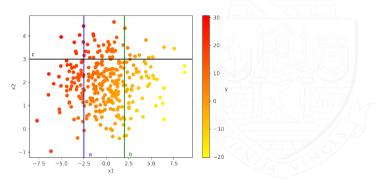
 - *x_j* ≤ *v* (left branch) *x_j* > *v* (right branch)





Decision splits

We need partitions that make the sub-spaces maximally homogeneous



How to pick the best threshold?



Variance reduction

- Regression trees usually aim at reducing the variance within the created partitions
- Variance Reduction heuristic

$$\mathsf{VR}(\mathsf{y},\mathsf{x},\Theta) = \mathsf{Var}(\mathsf{y}) - \frac{|\mathsf{y}_{\mathsf{x}} \leq \Theta|}{|\mathsf{y}|} \mathsf{Var}(\mathsf{y}_{\mathsf{x}} \leq \Theta) - \frac{|\mathsf{y}_{\mathsf{x}} > \Theta|}{|\mathsf{y}|} \mathsf{Var}(\mathsf{y}_{\mathsf{x}} > \Theta)$$

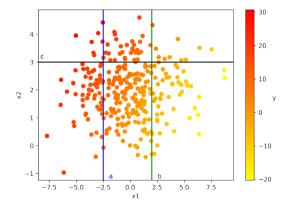
Best split candidate

$$(x_*, \Theta_*) = \operatorname{argmax}_{(x_i, v), x_i \in \{x_1, \dots, x_m\}, v \in \mathbb{R}} \mathsf{VR}(y, x_i, v)$$

- Equivalent to minimizing the Mean Squared Error (MSE)
 - The created partitions are maximally compact



Variance Reduction



►
$$VR(x_1, a) = 85.445 - (\frac{65}{300}) \times 22.895 - (\frac{235}{300} \times 59.75) = 33.683$$

► $VR(x_1, b) = 35.570$
► $VR(x_2, c) = 7.811$

We are all set to build regression trees!



Incremental regression trees: the needed tools

- 1. How to assure that our split candidates are indeed the best ones?
 - ► Hoeffding Bound √
- 2. How to evaluate split candidates?
 - 2.1 We need to calculate the elements of the VR equation
 - Incremental variance calculation!
 - 2.2 For any given partition (x_i, v) : $y_{x_i \le v}$ and $y_{x_i > v}$
 - How to do that incrementally and with reduced memory footprint? (and running time)
 - Answer: attribute observer (AO) algorithms (a.k.a. splitters)



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The naive approach

- ► Keep:
 - n: number of observations
 - $\sum y$: sum of y values
 - $\sum y^2$: sum of the squared y values
- $Var = \frac{1}{n-1} \left(\sum y^2 \frac{1}{n} \left(\sum x \right)^2 \right)$
- Used in Fast Incremental Model Tree with Drift Detection¹ (FIMT-DD)





Naive approach: the cool part

- Imagine that we keep two variance estimators:
 - Total variance of $y: var_T(y)$
 - ► Variance of *y* for elements that satisfy $x_i \le v$: $var_i(y)$ (left tree branch)
- How do we get the complement, $var_r(y)$? (right tree branch)
 - $n_r = n_T n_l$

$$\sum y_r = \sum y_T - \sum y_I$$

- $\sum y_r^2 = \sum y_T^2 \sum y_l^2$
- ▶ We do not need to keep *var_r*!
 - Memory savings to the attribute observers



Why naive?

- Both $\sum y^2$ and $\sum y$ can become really big
- Numerical cancellation
- Sometimes can even yield negative variance values (??)
- Text books do not advice to use this estimator in real-world applications¹

¹ Knuth, D.E., 2014. Art of computer programming, volume 2: Seminumerical algorithms. Addison-Wesley Professional.



Welford's algorithm: a stable solution

Initialize: $\overline{x}_1 = 0$, $M_{2,1} = 0$. For any n > 1:

$$\blacktriangleright \overline{x}_n = \overline{x}_{n-1} + \frac{x_n - \overline{x}_{n-1}}{n}$$

$$\blacktriangleright M_{2,n} = M_{2,n-1} + (x_n - \overline{x}_{n-1})(x_n - \overline{x}_n)$$

• Variance: $\frac{M_{2,n}}{n-1}$

In the next slide we drop the *n* indexing, for simplicity





Handling partial statistics

Handling addition¹ and subtraction²

Addition:

 $n_{AB} = n_A + n_B$ $\overline{x}_{AB} = \frac{n_A \overline{x}_A + n_B \overline{x}_B}{n_{AB}}$

$$\bullet M_{2,AB} = M_{2,A} + M_{2,B} + \delta^2 \frac{n_A n_B}{n_{AB}}$$

• In the expressions above, $\delta = \overline{x}_B - \overline{x}_A$

Subtraction:

 \blacktriangleright $n_A = n_{AB} - n_B$

 $\blacktriangleright \overline{x}_A = \frac{n_{AB}\overline{x}_{AB} - n_B\overline{x}_B}{n_A}$

• $M_{2,A} = M_{2,AB} - M_{2,B} - \delta^2 \frac{n_A n_B}{n_B}$

² Mastelini, S. M., de Leon Ferreira, A. C. P., 2021. Using dynamical quantization to perform split attempts in online tree regressors. Pattern Recognition Letters, 145, (pp. 37-42).



 n_{AB}

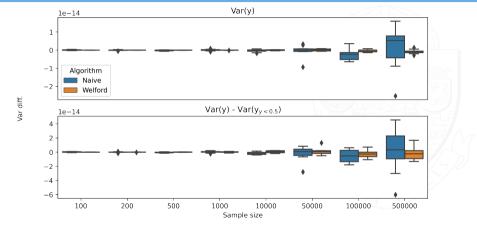
¹ Chan, T.F., Golub, G.H. and LeVeque, R.J., 1982. Updating formulae and a pairwise algorithm for computing sample variances. In COMPSTAT 1982 5th Symposium held at Toulouse 1982 (pp. 30-41). Physica, Heidelberg.

Benchmarking the variance estimators

- Naive and Welford against the non-incremental variance estimator
- Increasing sample size
- Difference between the obtained variances
 - Ground truth: non-incremental variance



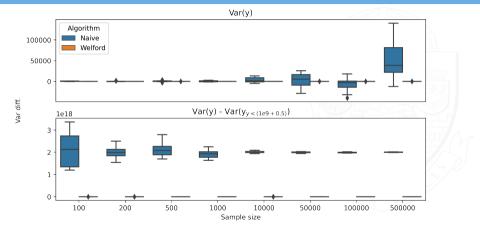
Benchmarks: uniform data between (0, 1)



► So far, so good



Benchmarks: adding constant shift of 10⁹



Hence, the Welford's algorithm will be our preferred choice

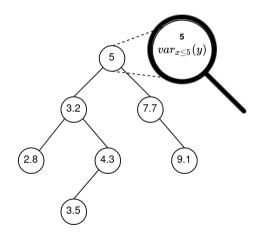


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E-BST: using trees to build trees

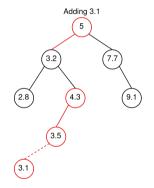




- Keep numerical features in a binary search tree:
 - Each node carries the feature value and a var estimator for y
- BST is not balanced



E-BST: adding a new node (I)

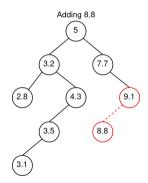




- The variance estimates are updated as new values are sorted down the BST
- Only statistics of the left branches are updated as new values are inserted into the BST
 - Partial statistics are kept
- The updated nodes are in red



E-BST: adding a new node (II)



The complete statistics are retrieved by performing a complete in-order traversal:

- 1. Create an auxiliary variance estimator varaux
- 2. If traversing to:
 - 2.1 left branch: pass varaux without modification
 - 2.2 **right** branch: update *var_{aux}* with the current node's statistics before descending
- The complete statistics (test ≤) are given by aggregating *var_{aux}* and the current node's variance estimator
- 4. Undo changes to varaux when backtracking



E-BST: cost and variant

	Insertion	$O(\log n)$ or $O(n)^*$
	Memory	<i>O</i> (<i>n</i>)
	Query time	<i>O</i> (<i>n</i>)
ne w	orst case, when i	incoming instances are ordered

Some alternatives to alleviate these costs:

* th

- Limit the number of nodes x
- ► Round the incoming data before insertion (Truncated E-BST TE-BST) √
- ► From time to time, remove bad split candidates from the BST¹ √

¹ Ikonomovska, E., Gama, J. and Džeroski, S., 2011. Learning model trees from evolving data streams. *Data mining and knowledge discovery*, 23(1), pp.128-168.



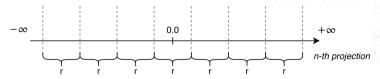
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QO: simple yet effective solution

- Can we get rid of the logarithmic cost per insertion?
 - ▶ What if we reached a cost of *O*(1) per insertion? Answer: Hashing!
- ► Inspiration in Locality Sensitive Hashing (LSH)¹:
 - Instead of mapping each element to its own hash slot, map similar elements to the same slot
- Straightforward projection rule: $h = \lfloor \frac{x_i}{r} \rfloor$, where *r* is the quantization radius (hyperparameter)



¹ Datar, M., Immorlica, N., Indyk, P. and Mirrokni, V.S., 2004, June. Locality-sensitive hashing scheme based on p-stable distributions. In Proceedings of the twentieth annual symposium on Computational geometry (pp. 253-262).



QO: example

Let's assume r = 0.25 and an empty hash table H

Insertion points: 2.3, 3.1, 7.78, 7.8

•
$$h_{2.3} = \lfloor \frac{2.3}{0.25} \rfloor = \lfloor 9.2 \rfloor = 9$$

•
$$h_{3.1} = \left[\frac{3.1}{0.25}\right] = 12$$

•
$$h_{7.78} = 31$$

•
$$h_{7.8} = 31$$



- ► For each slot we keep the mean *x* value and a variance estimator for *y*
 - Split points: Middle point between two consecutive slots



QO: cost and limitations

Cost	E-BST	TE-BST	QO
Insertion (per instance)	$O(\log n)$ or $O(n)^*$	$O(\log n')$ or $O(n')^*$	<i>O</i> (1)
Memory	<i>O</i> (<i>n</i>)	<i>O</i> (<i>n</i> [′])	<i>O</i> (<i>H</i>)
Query time	<i>O</i> (<i>n</i>)	<i>O</i> (<i>n</i> [′])	$O(H \log H)$

* the worst case, when incoming instances are ordered

- |H|: number of slots in the hash
- $n' \leq n$ (depending on the rounding procedure)
- The costs of QO depend on the choice of r



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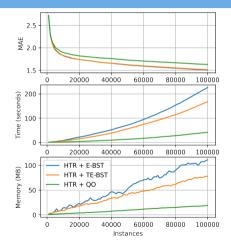


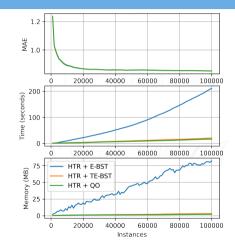
Friendly reminder

- Regression trees maximize the VR when making splits
- Incremental decision trees use attribute observers (AO) to evaluate split candidates
- Each tree leaf carries one AO per input feature



100k samples: Friedman and Planes2D







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- We apply a robust and incremental variance estimator to Hoeffding Tree Regressors
- We proposed a simple yet effective attribute observer algorithm for incremental regression tree construction
- QO is able to deliver faster tree construction with reduced memory footprint, while keeping the error increase minimal
- You can check everything in River :-D



What next?

- Mini-batches and distributed processing
 - QO is mergeable!
 - We can take advantage of vectorization and distributed processing units
- Investigate in depth the impact of r (quantization radius) and alternatives to select it automatically



Acknowledgments

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Thank you so much for your attention!

Questions?



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